





## The **BUB-Tree**

(bounding UB-Tree)



dealing with dead space

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## Motivation

Objective: Multidimensional indexing

- real data is always skewed
  - data warehousing
  - spatial data
- $\Rightarrow$  there is a lot of dead space
- UB-Tree partitions the whole universe
- ⇒ good performance in general, but queries on dead space suffer
- has the R-Tree a "better" partitioning?



use Z-intervals instead of Z-separators

- store (zstart,zend) in a B-Tree
- ⇒ higher index levels reflect Z-intervals of lower ones



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- split only when Z-gap volume exceeds a given % of Z-region volume
- split only when Z-gap volume exceeds a given volume



• we need two algorithms developed for UB-Trees

- NextJumpIn(z): calculates next intersection point greater than z of the Z-curve with the query box
- NextJumpOut(z): calculates the point greater than z (where Z is within the query box) on the Z-curve where it leaves the query box again

## Summary

- "twice" the space requirements for index part
  - reduced index fanout,
  - + but high potential for compression due to prefixes
- + "only" populated parts of the universe are indexed
- + UB-Tree basic algorithms can be reused
- + still a disjoint space partitioning
- + logarithmic cost for basic operations, i.e. it is a dynamic index structure which is not true for the R\*-Tree!
- + query performance equal or better to R\*-Tree
  - + also with dual space approach for GIS data