# D2P: Distance-Based Differential Privacy in Recommenders

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# ABSTRACT

The upsurge in the number of web users over the last two decades has resulted in a significant growth of online information. This information growth calls for recommenders that personalize the information proposed to each individual user. Nevertheless, personalization also opens major privacy concerns.

This paper presents D2P, a novel protocol that ensures a strong form of differential privacy, which we call distance-based differential privacy, and which is particularly well suited to recommenders.

D2P avoids revealing exact user profiles by creating *altered* profiles where each item is replaced with another one at some *distance*. We evaluate D2P analytically and experimentally on MovieLens and Jester datasets and compare it with other private and non-private recommenders.

# **1. INTRODUCTION**

In the modern generation of web-based services, the number of users is increasing exponentially. This number bumped up from 16 million users in 1995 to 3 billion users in 2014. The web has become a big storehouse of information (about 2.5 billion GB of data are created everyday), making it impossible for an individual to explore the whole web contents to extract relevant data. This clearly calls for *personalization* [4]. Personalizing the web led in turn to the advent of *recommenders* [17]. These systems filter out user-specific information in real-time, leveraging user activities and behaviors. Recommenders are primarily used in the context of e-commerce to suggest books, DVDs (*Amazon.com*), trips (*TripAdvisor*), music (*last.fm*) and even research papers (*Mendeley*). They are also used to filter user-specific news (*Google News, Yahoo News*).

However, the tendency towards personalization has raised a privacy concern [27] as more and more personal data is being collected and used. It is often observed that when an Internet user accesses some service, the provider of this service typically claims the ownership of any personal information provided by the user.

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The service provider sometimes even distributes the collected information to third parties like advertising and promotional partners [1]. Even the sharing of anonymised user information like the Netflix Prize dataset might end up not being secure. For instance, Narayanan et. al presented a de-anonymization attack that linked the records in the Netflix Prize dataset with the IMDB profiles available publicly [24].

Particularly fragile among recommenders are Collaborative Filtering (CF) ones [31]: these are widespread because of their ability to provide *serendipitous* recommendations (*unexpected* but *desired* recommendations) [27]. CF recommenders make predictions about the preferences of the users by collecting suggestions from similar users (user-based) or finding similar items (item-based) based on *neighborhood*. CF recommenders provide significant advantages over alternatives (e.g. content-based approaches [33]) due to the substantial number of features they use. Yet, CF recommenders are particularly vulnerable to privacy attacks as they rely on direct information about user *profiles* to provide good recommendations. They aggregate *user preferences* [28] in ways analogous to database queries, which can be exploited by adversaries to extract personal identifiable information about a specific user [27].

Clearly, CF recommenders induce an inherent tradeoff between privacy and quality [19]. In this paper, we address this tradeoff by exploring a promising approach where the information used for computing recommendations is *concealed*. We present D2P, a novel protocol that uses a probabilistic substitution technique to create the AlterEgo profile of an original user profile. D2P ensures a strong form of differential privacy [7, 8], which we call Distancebased Differential Privacy. Differential privacy [7, 8] is a celebrated property, originally introduced in the context of databases. Intuitively, it ensures that the removal of a record from a database does not change the result of a query to that database - modulo some arbitrarily small value ( $\epsilon$ ). In this sense, the presence in the database of every single record - possibly revealing some information about some user - is anonymous as no query can reveal the very existence of that record to any other user (modulo  $\epsilon$ ). Applying this notion in the context of recommenders would mean that - modulo  $\epsilon$  - no user Y would be able to guess - based on the recommendations she gets - whether some other user X has some item I in her profile, e.g., whether X has seen some movie I. Such a guarantee, however, might be considered too weak as nothing would prevent Y from guessing that X has in her profile some item that is very similar to I, e.g., that X has seen some movie similar to I.

We strengthen the notion of differential privacy in the context of CF recommenders to guarantee that any user Y is not only pre-

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vented from guessing whether the profile of X contains some item I, but also whether the profile of X contains any item I' within some *distance*  $\lambda$  from I (say any movie of the same category of I): hence the name **D***istance-based* **D***ifferential* **P***ivacy* (D2P). We present a protocol, named D2P, which ensures this property.

The basic idea underlying D2P is the following. We build, for each user profile, an *AlterEgo* profile corresponding to it. The latter profile is based on the former one where we probabilistically replace some of the items with either related or random ones. This poses of course a challenging technical problem. If the *AlterEgo* profile is too far from the original one, the recommendation quality is impacted: we lose the benefits of collaborative filtering. If the profile is too close to the original one, privacy remains weak.

We demonstrate in the paper that the quality of the D2P recommendation is still good for values of  $\lambda$  that can hide items within a reasonable distance from the original profile - what might be considered a reasonable distance depends on the dataset as we explain later in the paper.

To illustrate the basic idea, consider traces from Movielens and the scenario of Figure 1, with a total of 5 movies. Consider a user who likes Shawshank Redemption (SR). We compute the distance between the other 4 movies from SR based on their similarity (as shown later in Equation 1 in Section 2.1). D2P selects movies (for replacement) with distance less than the upper bound ( $\lambda = 0$ , 1 or 2) with high probability (p) and any random movie from the dataset, including those close to the item to be replaced, with a low probability (1 - p). If  $\lambda$  is set to 0, then D2P satisfies the classical differential privacy (with  $\epsilon$  given in Equation 12 in Section 3.2.3). Our results in Section 4 show that even if we consider  $\lambda$  as 6.5, we still have a good recommendation quality.

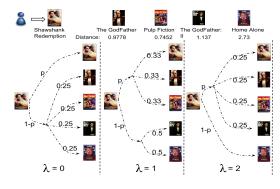


Figure 1: D2P Illustration.

We perform a thorough evaluation of  $D2\mathcal{P}$ : (i) we first analytically compute the guarantees ensured by  $D2\mathcal{P}$ , in terms of parameters  $\epsilon$  and  $\lambda$ , and then (ii) we evaluate experimentally the quality of recommendations provided on real datasets, namely Movie-Lens and Jester. Our results show that  $D2\mathcal{P}$  provides proved privacy guarantees while preserving the quality of the recommendation. We demonstrate, for instance, that  $D2\mathcal{P}$  achieves 1.5 times the coverage [10] provided by a standard recommender for Movie-lens dataset. Additionally, we show that the privatization overhead in  $D2\mathcal{P}$  is very small compared to [21], which makes it appealing for real-time workloads.

Interestingly, D2P is a generic protocol. As we show through our performance results, it applies well in the context of a userbased as well as an item-based recommender. D2P can also be customized for recommendation infrastructures where a *KNN* computation is deployed either on the cloud [26] or on user machines [5]. The rest of the paper is organized as follows. Section 2 presents D2P. Section 3 discusses its privacy guarantees. Section 4 evaluates the quality and coverage of its recommendations along with the privatization overhead. Section 5 discusses related work. We conclude the paper in Section 6 by discussing future works.

# **2.** D2P **RECOMMENDER**

D2P considers a general CF recommendation scheme based on KNN (K Nearest Neighbors [31]). The working principle of such a scheme is twofold. Firstly, the *k* most similar *neighbors* of any active user are identified in the KNN selection phase. Secondly, the recommendation algorithm is run to suggest items to the users leveraging the profiles obtained through the KNN selection.

### 2.1 Underlying Scheme

We consider a recommender scheme that stores *user profiles* and *item profiles*. The profile of a user  $\mathcal{U}$ , denoted by  $P_{\mathcal{U}}$ , consists of all the items rated (alternatively shared or liked) by  $\mathcal{U}$  along with the ratings. In our implementation, we convert the numerical ratings into binary ratings, a *like* (1) or a *dislike* (0). <sup>1</sup> An *item profile* ( $P_{\mathcal{I}}$ ) consists of users who rated item  $\mathcal{I}$  along with the ratings.

D2P relies on the *distance* between items to create *AlterEgo* profiles, as we discuss below. The recommender in D2P operates in four phases as shown in Figure 2.

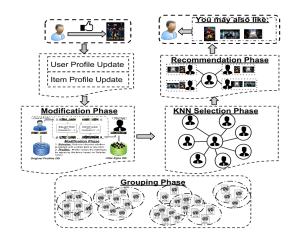


Figure 2: D2P Recommendation Scheme.

### 2.1.1 Grouping Phase.

In this phase, groups are formed for each item: group  $G_i$  for item i contains all the items with distance less than a predefined upperbound  $\lambda$ . In our scheme, we define the distance  $\Lambda_{i,j}$  between items i and j as:

$$\Lambda_{i,j} = \frac{1}{\Psi(i,j)} - 1 \tag{1}$$

Here,  $\Psi(i, j)$  denotes the cosine similarity between items *i* and *j*. The neighboring group  $G_j$  of a group  $G_i$  is defined as a group with which group  $G_i$  shares at least one item. Groups can also be formed based on item features (e.g. genres, date-of-release in case of movies) where similarity is measured between the feature vectors of the items. The groups need to be updated periodically to account for newly added items and ratings. In D2P, the grouping of

<sup>&</sup>lt;sup>1</sup>Binary ratings are considered for the sake of simplicity: this scheme can be generalized to numerical ratings.

the items in the *Grouping Phase* is performed by the *FormGroups* function shown in Algorithm 1. An item can be included in more than one groups, e.g., an *action-comedy* movie *X* can be present in the group of an *action* movie as well as in the group of a *comedy* movie.

**Algorithm 1 Grouping**: FormGroups(ItemSet): Grouping Phase where **ItemSet** is the set of all items in Database

1:	Parameter: $\lambda$	$\triangleright$ Distance threshold
2:	: var $ItemSet$ ; $\triangleright$ $Denotes set of$	all items in the network
3:	: var $\lambda$ ;	$\triangleright$ Distance Metric
4:	: for all <i>iid</i> : item in <i>ItemSet</i> do	
5:	: $Group_{iid}.add(iid);$	
6:	for all $rid$ : item in $(ItemSet \setminus i)$	id) do
7:	: $S = \Psi(iid, rid);$	$\triangleright$ Compute Similarity
8:	: if $S > 0$ then	
9:		
10:	: if $\Lambda_{iid,rid} \leq \lambda$ then	
11:	: $Group_{iid}.add(rid);$	
12:	end if	
13:	: end if	
14:	: end for	
15:	: end for	
16:	: return: $Group$ ; $\triangleright$	The groups for the items

### 2.1.2 Modification Phase.

D2P relies on the above mentioned groups of items to create *AlterEgo* profiles, avoiding to reveal the exact ones. The two core components of D2P are the *Selector*, which selects the items to replace and the *Profiler*, which determines by which items those entries should be replaced. The *AlterEgo* profile of a user U denotes the imitation profile of U which hides the user preferences by substituting items in the user profile by utilizing D2P. The *selector* and *profiler* are in charge of computing these *AlterEgo* profiles to preserve  $(\epsilon, \lambda)$ -differential privacy. Details about *selector* and *profiler* are provided later.

#### 2.1.3 KNN Selection Phase.

In user-based CF recommenders, a K-Nearest Neighbors (KNN) [31] algorithm computes the K most similar users based on some similarity metric. In this phase, we periodically update the top  $K_{users}$  similar users for an active user as the *neighbors* using the *AlterEgo* profiles generated in the modification phase.

#### 2.1.4 Recommendation Phase.

In this final phase, the recommendations are computed using those  $K_{users}$  neighbors. In the context of this paper, we select the most popular items among the *neighbors* of  $\mathcal{U}$  to be recommended to  $\mathcal{U}$ .

Some maintenance operations are needed: (i) *Profile update:* When a user  $\mathcal{U}$  rates an item  $\mathcal{I}$ , then both  $P_{\mathcal{U}}$  and  $P_{\mathcal{I}}$  are updated. Profiles are updated incrementally as in standard online recommenders. (ii) *Group Update:* The static nature of the relationship (similarity) [18, 29] between items stabilizes the grouping phase. So, the frequency of group updates has little impact on the quality of the provided recommendations; The groups are updated periodically after every 10 days in our evaluation. (iii) *Recommendation:* The new recommendations are delivered to the active user incrementally whenever an item is rated by the user. In  $\mathcal{D}2\mathcal{P}$ , only the *AlterEgo* profiles of the KNN are updated during each recommendation. We take into account the recent ratings provided by the users to compute recommendations. Privacy breaches occur in a standard user-based CF recommender due to leakage of the information of neighboring profiles to the active user through recommendations provided to her. D2P protects the privacy of users in the modification phase through two components: *Selector* and *Profiler* as conveyed by Figure 3. These two components conceal the neighbors' information from the active user, preventing this user to correlate the recommendations to the neighbors' profiles. The *selector* and *profiler* are responsible for forming the *AlterEgo* profiles in such a way that the quality is not impacted too much while privacy is preserved. We now provide details on these two core components.

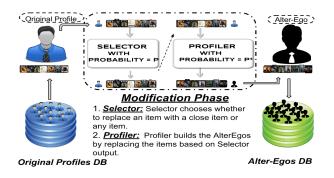


Figure 3: D2P Modification Phase.

## **2.2** D2P Components

### 2.2.1 D2P Selector

The *selector* is responsible for selecting the items to replace by the *profiler* to form the *AlterEgo* profiles. We select an item with a probability p to replace with any possible item at random and with a probability 1 - p to replace with some random item from the respective group (and neighboring groups) for that respective item. The *getSelectProb* function mentioned in Algorithms 2 and 3, returns a random real number between 0 and 1. Finally, the *selector* outputs a set of actual items (*GItems*) to be replaced by *GroupItems* and another set of actual items (*RItems*) to be replaced by any item from the set of all possible items at random.

**Algorithm 2 Selector Algorithm:** Selector  $(P_u)$  where  $P_u$  is the profile of user u

1: Parameter: p	$\triangleright$ Selector Probability
2: var $GItems[U] = NULL;$	$\triangleright$ Replace with group item
3: var $RItems[U] = NULL;$	
4: for all $iid$ : item in $P_U.getI$	tems() do
5: <b>if</b> $getSelectProb() > p$ <b>t</b>	hen
6: $GItems[U] = GItem$	$s[U] \cup iid;$
7: end if	
8: <b>if</b> $getSelectProb() \le p$ <b>t</b>	
9: $RItems[U] = RItem$	$s[U] \cup iid;$
10: <b>end if</b>	
11: end for	
12: return: $\{GItems[U], RItems[U], RItems[$	$ems[U]\};$

### 2.2.2 D2P Profiler

The *profiler* builds the *AlterEgo* profiles which are used in the KNN selection phase. The *profiler* replaces items in *Gltems* with items from their respective group (and neighboring groups) with a

probability  $1 - p^*$  and retains the original item with a probability  $p^*$ . We also substitute items in *RItems* with items from the set of all possible items with a probability  $1 - p^*$  and preserves the actual ones with a probability  $p^*$ . The *SRSI* (Select Random Set Item) function in Algorithm 3 selects randomly an item from the respective groups' items. It selects either from *GroupItems* (based on a distance metric between items) for all the items in the set *GItems* or from the *ItemSet* for all the items in *RItems*. In the following sections, we show that D2P ensures users' privacy while preserving a good recommendation quality.

**Algorithm 3 Profiler Algorithm**:  $Profiler(P_u)$  where  $P_u$  is the profile of user u

 $\triangleright$  Profiler Probability 1: Parameter:  $p^*$ 2: var {GItems[U], RItems[U]} =  $Selector(P_U)$ ; 3: var  $Items[U] = GPI(P_U);$  $\triangleright$  Get items from  $P_U$ var ItemSet;  $\triangleright$  Set of all items in the network 4: for all *iid* : item in  $P_U.getItems()$  do 5:  $GroupID = Group_{iid};$ 6:  $NBGroupIDs = \overline{Group_{iid}}.getNeighbors();$ 7: 8:  $Groups = GroupID \cup NBGroupIDs;$  $\begin{array}{l} Group I = Group I D S, \\ Group I tems = \bigcup_{G \in Groups} Group.get(G); \\ \text{if } (getSelectProb() > p^* \& iid \in GItems[U]) \text{ then} \\ rid = SRSI(iid, Group I tems); \end{array}$ 9: 10: 11: end if 12: if  $(getSelectProb() > p^* \& iid \in RItems[U])$  then 13:  $\ddot{rid} = SRSI(iid, ItemSet);$ 14: 15: end if  $P_U = (P_U \setminus iid) \cup rid;$ 16: 17: end for 18: return:  $P_U$ ;  $\triangleright$  AlterEgo profile for user U

Interestingly, D2P can also be applied in recommendation infrastructures where the KNN is computed by third-party cloud services that act as intermediaries between the recommendation server and users: these servers create the *AlterEgo* profiles, preserving privacy with respect to a server. Moreover, D2P can be applied by the users themselves (in P2P or hybrid infrastructures [5]), preserving privacy of users against other users.

# 3. PRIVACY GUARANTEES

Preserving privacy in CF recommenders is challenging. It was shown using the Netflix Prize dataset that even anonymizing individual data before releasing it publicly is not enough to preserve privacy [24]. Even cryptographic approaches do not preclude the possibility of the output leaking information about the personal input of individuals [32]. The need for stronger and robust privacy guarantees motivated the emergence of the notion of *Differential Privacy* [7, 8, 9]. First introduced in the context of databases, differential privacy provides quantifiable privacy guarantees. We introduce a stronger form of this notion in the context of recommenders by accounting for the concept of distance between items.

# **3.1** Privacy for Recommenders

### 3.1.1 Differential Privacy

Differential Privacy (DP) implies that the output of a given function becomes significantly more or less likely - based on some parameter  $\epsilon$  - if the inputs differ in one record. The basic intuition is that an observer can extract limited information from the output in the absence or presence of a specific record in the database. DEFINITION 1. (DIFFERENTIAL PRIVACY) A randomized function  $\mathcal{R}$  provides  $\epsilon$ -differential privacy if for all datasets  $\mathcal{D}_1$  and  $\mathcal{D}_2$ , differing on at most one element, and all  $S \subseteq Range(\mathcal{R})$ , the following inequality always holds:

$$\frac{Pr[\mathcal{R}(\mathcal{D}_1) \in \mathcal{S}]}{Pr[\mathcal{R}(\mathcal{D}_2) \in \mathcal{S}]} \le e^{\epsilon}$$

Here,  $e^{\epsilon}$  denotes  $exp(\epsilon)$ .

#### 3.1.2 Distance-based Differential Privacy

With differential privacy applied in its classical form recalled above to a recommender, an adversary (a curious user) cannot know if one item has been rated by a user. However, the adversary can know about items similar to the rated ones. Hence, the adversary can infer fairly accurate information about user preferences without knowing the exact items rated by that user. In this sense, classical differential privacy is not enough in the context of a recommender.

Our notion of *Distance-based Differential Privacy* is stronger: it extends DP to recommenders. We ensure differential privacy for all the items, rated by that user, and ones that are within a distance of  $\lambda$ . The distance parameter ( $\lambda$ ) determines the *closely related items to form the AlterEgo profiles*, thereby *concealing* the actual user profiles and preferences. The distance parameter also aids in tuning the recommendation quality using the *AlterEgo* profiles as shown later in Figures 12a and 12b.

It is important to notice that our notion of Distance-based differential privacy is independent from the underlying recommendation algorithm used. To define this new notion more precisely, we first define the notions of *Distance-based Group* and *Adjacent Profile Sets*.

DEFINITION 2. (GROUP DEFINITION ELEMENT-WISE) We denote by  $\mathbb{E}$  the set of all elements. For every element  $x \in \mathbb{E}$ , distance function  $\Lambda : \mathbb{E} \times \mathbb{E} \to \mathbb{R}^+ \cup \{0\}$ , and fixed distance threshold  $\lambda$ , then  $\mathcal{GRP}_{\lambda}(x)$  is defined as the collection of all elements  $x_k \in \mathbb{E}$ such that  $\Lambda_{x,x_k} \leq \lambda$ . More specifically:

$$\mathcal{GRP}_{\lambda}(x) = \{x_k \in \mathbb{E} | \Lambda_{x,x_k} \le \lambda\}$$

We extend this notion of groups to a set of elements where each element in the set has a *Group* defined by Definition 2.

DEFINITION 3. (GROUP DEFINITION SET-WISE) For a set of elements S,  $\mathcal{GRP}_{\lambda}(S)$  is the union of all the groups:  $\mathcal{GRP}_{\lambda}(s)$  for each element  $s \in S$ . More specifically:

$$\mathcal{GRP}_{\lambda}(\mathcal{S}) = \bigcup_{s \in \mathcal{S}} \mathcal{GRP}_{\lambda}(s)$$

We now introduce the notion of *Neighboring Groups* (used in Section 3.2.3).

DEFINITION 4. (NEIGHBORING GROUP) We define the  $\mathcal{KNN}$  groups  $(\mathcal{KNN}(\mathcal{GRP}_{\lambda}(x)))$  of  $\mathcal{GRP}_{\lambda}(x)$  for an element x as the  $Top-\mathcal{K}$  groups sorted in decreasing order by the count of shared elements with  $\mathcal{GRP}_{\lambda}(x)$ .

DEFINITION 5. (ADJACENT PROFILE SET) An event in the context of D2P is an interaction between the system and the user when the user provides a rating for some item in the system. Two profile sets  $D_1$  and  $D_2$  as adjacent profile sets when  $D_1$  and  $D_2$  differ in only one event, which implies one user-item rating pattern is different in these two profile sets.

	Notations			
$\mathcal{M}$	A mechanism relying on AlterEgo profiles to provide recom-			
	mendations			
U	User to whom some recommendation is provided			
$\mathcal{P}_i$	Accurate $i^{th}$ profile from profile set $\mathcal{D}$			
$\mathcal{P}_i^r$	AlterEgo profile of original profile $\mathcal{P}_i$			
$p_i^r$	Any random possible $AlterEgo$ profile for original profile $\mathcal{P}_i$			
$\mathcal{D}$	Original Profile Set			
$\mathcal{D}'$	Profile Set which differs with $\mathcal{D}$ at an event			
S	Arbitrary set of elements			
$\mathcal{N}_{\mathbb{E}}$	Total number of elements			
π	Any possible permutation of numbers in the range $\{0,n\}$			

Table 1: Notations

For any arbitrary recommendation mechanism  $\mathcal{R}$ , which takes a profile set and a specific user as input, the output is the set of items that the algorithm recommends to that specific user.

DEFINITION 6. (DISTANCE-BASED DIFFERENTIAL PRIVACY) For any two adjacent profile sets  $D_1$  and  $D_2$ , where U denotes any arbitrary user and S denotes any possible subset of elements, then any mechanism  $\mathcal{R}$  is  $(\epsilon, \lambda)$ -private if the following inequality holds:

$$\frac{Pr[\mathcal{R}(\mathcal{D}_1,\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]}{Pr[\mathcal{R}(\mathcal{D}_2,\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]}\leq e^{\epsilon}$$
(2)

The result of the recommendations for two profile sets that are close to each other are of the same order probabilistically with a coefficient of  $e^{\epsilon}$ . Later in Section 3.2.3, we present the mathematical relationship between  $\epsilon$  and  $\lambda$ . D2P conceals the profiles by anonymizing elements within distance  $\lambda$  from the elements of the original profile. We get the classic notion of differential privacy with  $\lambda$  as 0. If we increase  $\lambda$  then the privacy increases but the quality decreases slightly as shown later in Figure 12a. In a user-level privacy scheme, more than one event can differ for a profile in two adjacent profile sets, whereas in an event-level privacy approach a single event differs for a profile in two adjacent profile sets. The proofs in the following section assume event-level privacy.

#### **3.2** Privacy Analysis

In this section, we analyze our D2P protocol based on the privacy parameter  $\epsilon$  and the distance parameter  $\lambda$ .

### 3.2.1 Notations

We fix an arbitrary user  $\mathcal{U}$  to whom we provide some recommendations.  $\mathcal{D}$  and  $\mathcal{D}'$  are two adjacent profile sets. Additional notations used later in the proofs are mentioned in Table 1.

#### 3.2.2 Proofs

As  $\mathcal{D}$  and  $\mathcal{D}'$  are two adjacent profile sets, using Definition 5 we know that  $\mathcal{D}$  and  $\mathcal{D}'$  differ at one event which is one user-item rating pattern in a profile  $\mathcal{P}_i$  denotes this profile in profile set  $\mathcal{D}$ whereas  $\mathcal{P}'_i$  denotes this profile in profile set  $\mathcal{D}'$ .  $\mathcal{P}_i$  has an element  $i_0$ , for which in  $\mathcal{P}'_i$  there is another element  $i'_0$  in place of  $i_0$ . So, exactly one rating pattern is different in  $\mathcal{P}_i$  and  $\mathcal{P}'_i$ . Let the elements in  $\mathcal{P}_i$  be  $i_0, i_1, ..., i_n$  and the elements in  $\mathcal{P}'_i$  be  $i'_0, i_1, ..., i_n$ .

**PROPOSITION** 1. For any given distance metric  $\lambda$  and any two elements *i* and *j*, we denote SUB(i, j) the event of substituting element *i* with *j* in a mechanism  $\mathcal{M}$ . This substitution probability is denoted by Pr(SUB(i, j)). Then, for mechanism  $\mathcal{M}$ , we propose

 $\epsilon$  as:

$$\begin{aligned} \epsilon &= \ln(\delta) \end{aligned} \\ \text{where } \delta &= \max_{i,j,k \in \mathbb{E} \text{ and } i \neq j} \left( \frac{\Pr(\mathcal{SUB}(i,k))}{\Pr(\mathcal{SUB}(j,k))} \right) \end{aligned}$$

THEOREM 1. Any mechanism  $\mathcal{M}$ , that relies on the AlterEgo of a profile  $\mathcal{P}_i$ , is an  $(\epsilon, \lambda)$ -private mechanism for  $\epsilon$  given in Proposition 1.

First, we present some lemmas and corollaries needed to prove Theorem 1.

LEMMA 1. For three arbitrary elements  $i, j, k \in \mathbb{E}$  and  $i \neq j$ , we get the following inequality:

$$\frac{Pr(\mathcal{SUB}(i,k))}{Pr(\mathcal{SUB}(j,k))} \le e^{\epsilon}$$

PROOF. Since, any ratio is upper-bounded by its maximum. Hence, from Proposition 1 we have:

$$\frac{Pr(\mathcal{SUB}(i,k))}{Pr(\mathcal{SUB}(j,k))} \le \left(\max_{i,j,k \in \mathbb{E} \text{ and } i \neq j} \frac{Pr(\mathcal{SUB}(i,k))}{Pr(\mathcal{SUB}(j,k))}\right) = e^{i k t}$$

Therefore, we get:

$$\frac{Pr(\mathcal{SUB}(i,k))}{Pr(\mathcal{SUB}(j,k))} \le e^{\epsilon}$$

LEMMA 2. Let  $\mathcal{M}$  be a privacy preserving mechanism which creates the AlterEgo  $(\mathcal{P}_i^r)$  of a profile  $\mathcal{P}_i$ .  $Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r))$  denotes the probability of the substitution event  $(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r))$  to create the AlterEgo. Then, the following inequality holds:

$$\frac{Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r))}{Pr(\mathcal{PS}(\mathcal{P}'_i, \mathcal{P}^r_i))} \le e^{\epsilon}$$

PROOF. Denote by  $\pi$  any permutation of  $\{0,...,n\}$ . Let  $P_{i_{\pi}}^{r}$  be a profile with elements  $i_{\pi(0)}^{r}$ ,  $i_{\pi(1)}^{r}$ , ...,  $i_{\pi(n)}^{r}$ .  $P_{i_{\pi}}^{r}$  denotes one possible permutation of the elements of  $\mathcal{P}_{i}^{r}$ . So, for any permutation  $\pi$ , we compute the probability of changing  $\mathcal{P}_{i}$  to  $\mathcal{P}_{i_{\pi}}^{r}$ . Then, we sum over all possible permutations to get the final probability of the substitution event ( $\mathcal{PS}(\mathcal{P}_{i}, \mathcal{P}_{i}^{r})$ ).

$$Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r)) = \sum_{\forall \pi} Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_{i_{\pi}}^r))$$
(3)

Now, based on the fact that every element is replaced independently of the replacement of other elements, we compute:

$$Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_{i_{\pi}}^r)) = \prod_{k=0}^n Pr(\mathcal{SUB}(i_k, i_{\pi(k)}^r))$$
(4)

Based on Equations 3 and 4, we get:

$$Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r)) = \sum_{\forall \pi} \left( \prod_{k=0}^n Pr(\mathcal{SUB}(i_k, i_{\pi(k)}^r)) \right)$$
(5)

In the same way, for  $\mathcal{P}'_i$ , we have:

$$Pr(\mathcal{PS}(\mathcal{P}'_{i}, \mathcal{P}'_{i})) = \sum_{\forall \pi} Pr(\mathcal{SUB}(i'_{0}, i^{r}_{\pi(0)})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i^{r}_{\pi(k)})) \quad (6)$$

Now, from Equations 5 and 6, we get:

$$\frac{Pr(\mathcal{PS}(\mathcal{P}_{i}, \mathcal{P}_{i}^{r}))}{Pr(\mathcal{PS}(\mathcal{P}_{i}^{\prime}, \mathcal{P}_{i}^{r}))} = \frac{\sum_{\forall \pi} Pr(\mathcal{SUB}(i_{0}, i_{\pi(0)}^{r})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i_{\pi(k)}^{r})))}{\sum_{\forall \pi} Pr(\mathcal{SUB}(i_{0}^{\prime}, i_{\pi(0)}^{r})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i_{\pi(k)}^{r})))} \quad (7)$$

From Lemma 1, we get for any three arbitrary elements  $i_0$ ,  $i'_0$  and  $i^r_0$ :

$$\frac{Pr(\mathcal{SUB}(i_0, i_0^r))}{Pr(\mathcal{SUB}(i'_0, i_0^r))} \le e^{\epsilon}$$
(8)

So, from Equations 7 and 8, we get:

$$\frac{Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r))}{Pr(\mathcal{PS}(\mathcal{P}'_i, \mathcal{P}_i^r))}$$

$$\leq \frac{\sum_{\forall \pi} e^{\epsilon} \cdot Pr(\mathcal{SUB}(i'_{0}, i^{r}_{\pi(0)})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i^{r}_{\pi(k)}))}{\sum_{\forall \pi} Pr(\mathcal{SUB}(i'_{0}, i^{r}_{\pi(0)})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i^{r}_{\pi(k)}))}$$
$$= e^{\epsilon} \cdot \frac{\sum_{\forall \pi} Pr(\mathcal{SUB}(i'_{0}, i^{r}_{\pi(0)})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i^{r}_{\pi(k)}))}{\sum_{\forall \pi} Pr(\mathcal{SUB}(i'_{0}, i^{r}_{\pi(0)})) \cdot \prod_{k=1}^{n} Pr(\mathcal{SUB}(i_{k}, i^{r}_{\pi(k)}))}$$

 $= e^{\epsilon}$ 

Hence, we can conclude that:

$$\frac{Pr(\mathcal{PS}(\mathcal{P}_i, \mathcal{P}_i^r))}{Pr(\mathcal{PS}(\mathcal{P}_i', \mathcal{P}_i^r))} \leq e^{\epsilon}$$

COROLLARY 1. For any two adjacent profile sets  $\mathcal{D}, \mathcal{D}'$  and any arbitrary AlterEgo profile set  $\mathcal{D}^r$ , we denote by  $\mathcal{DS}(\mathcal{D}, \mathcal{D}^r)$  the Set Substitution event for  $\mathcal{D}$  with  $\mathcal{D}^r$ . We assume the profile sets  $\mathcal{D}$ and  $\mathcal{D}'$  differ at the *i*<sup>th</sup> profile (*P<sub>i</sub>*). Then, we have:

$$\frac{Pr(\mathcal{DS}(\mathcal{D},\mathcal{D}^r))}{Pr(\mathcal{DS}(\mathcal{D}',\mathcal{D}^r))} \leq e$$

PROOF. Using the same approach as in Lemma 2, we get:

$$\frac{Pr(\mathcal{DS}(\mathcal{D},\mathcal{D}^r))}{Pr(\mathcal{DS}(\mathcal{D}',\mathcal{D}^r))} = \frac{\sum\limits_{\forall \pi} Pr(\mathcal{PS}(P_i,P_{\pi(i)}^r)) \cdot \prod\limits_{k=1,k\neq i}^m Pr(\mathcal{PS}(P_k,P_{\pi(k)}^r))}{\sum\limits_{\forall \pi} Pr(\mathcal{PS}(P_i',P_{\pi(i)}^r)) \cdot \prod\limits_{k=1}^m Pr(\mathcal{PS}(P_k,P_{\pi(k)}^r))}$$

From Lemma 2, we have:

$$\frac{Pr(\mathcal{PS}(P_i, P_{\pi(i)}^r))}{Pr(\mathcal{PS}(P_i', P_{\pi(i)}^r))} \le e^{\epsilon}$$

Therefore, using Lemma 2, we get:

$$\frac{Pr(\mathcal{DS}(\mathcal{D},\mathcal{D}^r))}{Pr(\mathcal{DS}(\mathcal{D}',\mathcal{D}^r))} \leq \frac{\sum\limits_{\forall \pi} \mathbf{e}^{\epsilon} \cdot Pr(\mathcal{PS}(P_i',P_{\pi(i)}^r)) \cdot \prod\limits_{k=1,k\neq i}^m Pr(\mathcal{PS}(P_k,P_{\pi(k)}^r))}{\sum\limits_{\forall \pi} Pr(\mathcal{PS}(P_i',P_{\pi(i)}^r)) \cdot \prod\limits_{k=1,k\neq i}^m Pr(\mathcal{PS}(P_k,P_{\pi(k)}^r))}$$

Hence, we get:

$$\frac{Pr(\mathcal{DS}(\mathcal{D}, \mathcal{D}^r))}{Pr(\mathcal{DS}(\mathcal{D}', \mathcal{D}^r))} \le e^{\epsilon}$$

PROOF OF THEOREM 1. Let  $\mathcal{D}$  and  $\mathcal{D}'$  be any two adjacent profile sets and  $\mathcal{D}^r$  be any arbitrary *AlterEgo* profile set. We define  $\mathcal{M}'$ as any mechanism which takes as input the *AlterEgo* profiles and the target user to whom we provide recommendations as shown in Figure 4. So, we can rewrite:

$$Pr[\mathcal{M}(\mathcal{D},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})] = \sum_{\mathcal{D}^{r}} Pr(DS(\mathcal{D},\mathcal{D}^{r})).Pr[\mathcal{M}'(\mathcal{D}^{r},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]$$
(9)

Using the same approach for the profile set  $\mathcal{D}'$ , we get the following equation:

$$Pr[\mathcal{M}(\mathcal{D}',\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})] = \sum_{\mathcal{D}^{r}} Pr(DS(\mathcal{D}',\mathcal{D}^{r})).Pr[\mathcal{M}'(\mathcal{D}^{r},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})] \quad (10)$$

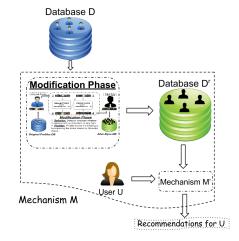


Figure 4: Relation between mechanisms M and M'.

From Equations 9 and 10, we arrive at:

$$\begin{split} & \frac{Pr[\mathcal{M}(\mathcal{D},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]}{Pr[\mathcal{M}(\mathcal{D}',\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]} \\ &= \frac{\sum_{\mathcal{D}^{T}} Pr(DS(\mathcal{D},\mathcal{D}^{r})).Pr[\mathcal{M}'(\mathcal{D}^{r},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]}{\sum_{\mathcal{D}^{T}} Pr(DS(\mathcal{D}',\mathcal{D}^{r})).Pr[\mathcal{M}'(\mathcal{D}^{r},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]} \end{split}$$

- Let  $\mathcal{Q}_{\mathcal{D}^r}$  denote the event probability:  $Pr[\mathcal{M}'(\mathcal{D}^r, \mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]$ , ) so we can reformulate the above equation as:

$$\frac{Pr[\mathcal{M}(\mathcal{D},\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]}{Pr[\mathcal{M}(\mathcal{D}',\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]} = \frac{\sum_{\mathcal{D}^{r}}Pr(DS(\mathcal{D},\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}}{\sum_{\mathcal{D}^{r}}Pr(DS(\mathcal{D}',\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}}$$
(11)

From Equation 11 and Corollary 1, we have:

$$\frac{\Pr[\mathcal{M}(\mathcal{D},\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]}{\Pr[\mathcal{M}(\mathcal{D}',\mathcal{U}) \in \mathcal{GRP}_{\lambda}(\mathcal{S})]} = \frac{\sum_{\mathcal{D}^{r}} \Pr(DS(\mathcal{D},\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}}{\sum_{\mathcal{D}^{r}} \Pr(DS(\mathcal{D}',\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}} \\ \leq \frac{\sum_{\mathcal{D}^{r}} e^{\epsilon}.\Pr(DS(\mathcal{D}',\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}}{\sum_{\mathcal{D}^{r}} \Pr(DS(\mathcal{D}',\mathcal{D}^{r})).\mathcal{Q}_{\mathcal{D}^{r}}} = e^{\epsilon}$$

Hence, using  $\epsilon$  from Proposition 1 for mechanism  $\mathcal{M}$ , we get:

$$\frac{Pr[\mathcal{M}(\mathcal{D},\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]}{Pr[\mathcal{M}(\mathcal{D}',\mathcal{U})\in\mathcal{GRP}_{\lambda}(\mathcal{S})]}\leq e^{\epsilon}$$

Therefore, using Definition 6, we can conclude that mechanism  $\mathcal{M}$  satisfies  $(\epsilon, \lambda)$ -privacy.  $\Box$ 

#### 3.2.3 Privacy Analysis of D2P

Now, we analyze our D2P privacy in the abstract model which we introduced in Section 2.

First, we denote the GroupItems for an item i in Algorithm 3 as:

$$\mathcal{G}_{\lambda}(i) = (\cup_{j \in \mathcal{KNN}(\mathcal{GRP}_{\lambda}(i))} \mathcal{GRP}_{\lambda}(j)) \cup \mathcal{GRP}_{\lambda}(i)$$

As mentioned in Section 2, the *selector* selects to replace an element *s* with any random element from  $\mathbb{E}$  with a probability *p* and with any random element from  $\mathcal{G}_{\lambda}(s)$  with a probability 1 - p. So, it finally outputs two sets of elements *GItems* and *RItems* for each user profile. For both of these sets (*GItems* and *RItems*), the *profiler* retains the original elements with probability  $p^*$ . It replaces elements in *GItems* with elements from  $\mathcal{G}_{\lambda}(s)$  and elements in *RItems* with any possible element  $e \in \mathbb{E}$  with probability  $1 - p^*$ . Here  $\mathcal{N}_{\mathbb{E}}$  is the total number of elements in  $\mathbb{E}$ .

So, we compute Pr(SUB(s, t)), for any two arbitrary elements *s* and *t*, in this abstract model. We get the following:

$$Pr(\mathcal{SUB}(s,t)) = \begin{cases} p^* + \frac{(1-p)(1-p^*)}{|\mathcal{G}_{\lambda}(s)|} + \frac{p(1-p^*)}{\mathcal{N}_{\mathbb{E}}} & \text{if } s = t\\ \frac{(1-p)(1-p^*)}{|\mathcal{G}_{\lambda}(s)|} + \frac{p(1-p^*)}{\mathcal{N}_{\mathbb{E}}} & \text{if } t \in \mathcal{G}_{\lambda}(s) \setminus s\\ \frac{p(1-p^*)}{\mathcal{N}_{\mathbb{E}}} & \text{if } t \notin \mathcal{G}_{\lambda}(s) \,. \end{cases}$$

Let  $\epsilon_{\mathcal{D2P}}^{(p,p^*,\lambda)}$  denote the  $\epsilon$  for  $\mathcal{D2P}$  with privacy parameters  $(p, p^* \text{ and } \lambda)$  and  $|\mathcal{G}_{\lambda}|$  denote  $\min_{s \in E} (|\mathcal{G}_{\lambda}(s)|)$ . Then, using the above substitution probabilities, we get:

$$\epsilon_{\mathcal{D}2\mathcal{P}}^{(p,p^*,\lambda)} = \ln(1 + \frac{p^* + \frac{(1-p)(1-p^*)}{|\mathcal{G}_{\lambda}|}}{\frac{p(1-p^*)}{\mathcal{N}_{\mathbb{R}}}})$$
(12)

So, when we compute using the original profile, we have  $p^* = 1$ , which implies  $\epsilon_{D2\mathcal{P}}^{(p,1,\lambda)} = \infty$  (*no privacy*). When  $p^* = 0$  in Equation 12, so all the items are replaced with some items. Then we have  $\epsilon_{D2\mathcal{P}}^{(p,0,\lambda)}$  as :

$$\epsilon_{\mathcal{D}2\mathcal{P}}^{(p,0,\lambda)} = \ln(1 + \frac{\frac{(1-p)}{|\mathcal{G}_{\lambda}|}}{\frac{p}{\mathcal{N}_{\mathbb{E}}}}) = \ln(1 + \frac{(1-p)\mathcal{N}_{\mathbb{E}}}{p.|\mathcal{G}_{\lambda}|})$$
(13)

From this  $\epsilon_{\mathcal{D}2\mathcal{P}}^{(p,0,\lambda)}$ , we see that when p increases, the probability to replace an item with a random item increases leading to more privacy and that is evident from the decreasing value of  $\epsilon_{\mathcal{D}2\mathcal{P}}^{(p,0,\lambda)}$  in Equation 13. When p = 1 in Equation 13,  $\mathcal{D}2\mathcal{P}$  achieves  $\epsilon_{\mathcal{D}2\mathcal{P}}^{(1,0,\lambda)} = 0$  (*perfect privacy*). For larger  $\lambda$ , the size of the groups becomes larger, hence privacy increases resulting in smaller  $\epsilon_{\mathcal{D}2\mathcal{P}}$ .

# 4. EXPERIMENTAL EVALUATION

This section presents an exhaustive experimental evaluation of our D2P protocol using two real datasets namely Jester and Movie-Lens. In particular, we compare the recommendation quality and coverage [10] of D2P with that of a non-private protocol directly relying on the original user profiles. We also provide a comparison with [21], one of the closest to our work. Additionally, we discuss an item-based version of D2P: *i*-D2P which we implemented and evaluated.

# 4.1 Experimental Setup

#### 4.1.1 Evaluation Metrics

We measure the recommendation quality as follows: we divide the dataset into a training set (80%) of the dataset trace) and a test set (20%). For each rating in the test set, a set of top recommendations is selected as the *Recommendation Set* (RS). We denote the size of

	Relevant	Irrelevant	Total
Recommended	tp	fp	tp + fp
Not Recommended	fn	tn	fn + tn
Total	tp + fn	fp + tn	N

Table 2: Confusion Matrix for true/false positive/negative recommendations.

the recommendation set as N. The quality is measured in terms of standard classification accuracy metrics (CAM) [30]. More precisely, we evaluate the extent to which the recommender is able to predict the content of the test set while having computed the KNN on the training set. We use *Precision* and *Recall* as classification accuracy metrics for they are conventionally used in top-N recommenders [6]. Table 2 shows the terms needed for defining *Precision* and *Recall: True Positives* (tp), *True Negatives* (tn), *False Positives* (fp), *False Negatives* (fn).

*Precision* or *True Positive Accuracy (TPA)* is the ratio of the number of relevant recommended items to the total number of recommended items.

$$Precision = TPA = \frac{tp}{tp+fp}$$

*Recall* or *True Positive Rate (TPR)* is the ratio of the number of relevant recommended items to the total number of relevant items.

$$Recall = TPR = \frac{tp}{tp+fn}$$

To get an estimate of the drop in quality, we measure the decrease in precision for Top-5 recommendations [22] (denoted by Pr@5), as most recommenders follow Top-N recommendations, e.g: *IMDB* uses *Top-6* list to suggest similar movies, *Amazon* uses *Top-4* list to suggest similar products and *last.fm* uses *Top-5* list to suggest similar music.

 $F_1$ -Score is used to access precision and recall simultaneously. Mathematically, it is the harmonic mean of *Precision* and *Recall*.

$$F_1 - Score = 2. \frac{Precision.Recall}{Precision+Recall}$$

### 4.1.2 Datasets

We evaluate D2P with two datasets: the MovieLens (ML) dataset [23] and the Jester one [14]. The ML dataset consists of 100, 000 (100K) ratings given by 943 users over 1682 movies. The Jester dataset [14] contains 4.1 million ratings of 100 jokes from 73,421 users. We use a subset of the Jester dataset with around 36K ratings given by 500 users over 100 jokes. The Jester subset consists of 500 users selected uniformly at random among all users who rated at least 50 jokes. D2P relies on the item-replacement technique, so the quality of the recommendation provided by D2P depends on how much two items are connected in the dataset. We thus consider datasets with diverse characteristics to evaluate D2P.

*Diversity:* We created 4 diverse datasets from the ML 100K dataset to cover a variety of characteristics (typically sparsity).

The ratings are stored in a user-item matrix where the rows of the matrix contain the user-ids and the columns contain the itemids. Then, the rows are sorted based on the total number of ratings given by the users and the columns are sorted based on the total number of times the items have been rated by different users. The partitioning of the dataset is shown in Figure 5 as *users*  $\times$  *items* matrix.

*Characterization.* To evaluate D2P in different settings, we characterize the datasets according to *Rating Density* metric. The Rating Density (RD) is the ratio of the number of ratings given by

	MLV	MLV
	mile 4	III.V
	영양 방법 문제 이 가지?	
	황료화관 것은 것 같아?	
· 建制行动的 1000 1000 1000 1000 1000	MLV <sub>2</sub>	MLA

Figure 5: ML1 Dataset Partitions based on rating density.

Dataset	#Users	#Items	Ratings	<b>RD</b> (%)
Jester	500	100	36000	71.01
ML1	940	1680	99647	6.31
$MLV_1$	470	840	76196	19.3
MLV <sub>2</sub>	470	840	16187	4.1
MLV <sub>3</sub>	470	840	6317	1.6
$MLV_4$	470	840	750	0.19

Table 3: Datasets characteristics.

the users in the dataset to the total number of ratings possibly given (number of users multiplied by the number of items).

Table 3 depicts the rating densities of different datasets.

## 4.2 Results

#### 4.2.1 Impact of Rating Density

Figure 6 shows the recall measured with varying size of the recommendation set in D2P with parameters p = 0.5,  $p^* = 0.5$  and  $\lambda = 1$ . We observe that higher rating density results in better recall using D2P. As shown in Table 3, the rating density of the Movielens 100K dataset is 6.31% and that of its 4 subsets varies with a maximum of 19.3% and minimum of about 0.19%. From Figure 6, we observe that D2P is not suitable for datasets with too low rating densities, like  $MLV_3$  and  $MLV_4$ , as these result in lower recall. However, we observe, for  $MLV_2$ , D2P provides slightly better recall compared to a more dense dataset (like  $MLV_1$ ). This happens because the number of items relevant to a user (in the test set) is less in  $MLV_2$  (more sparse) compared to  $MLV_1$  (less sparse). However, for more sparse datasets like  $MLV_3$  or  $MLV_4$ , collaborative filtering is not effective because the ratings are insufficient to identify similarities in user interests.

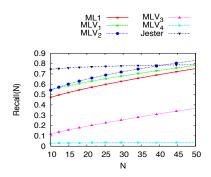


Figure 6: Recall@N with varying Dataset Characteristics.

### 4.2.2 Privacy-Quality Trade-off

#### *Effect of Profiler probability* $(p^*)$ *.*

We vary the value of parameter  $p^*$  from the *Profiler* algorithm from a minimum of 0 to a maximum of 1 (*no privacy*) with other parameters  $\lambda = 1, p = 0.5$ .

*Movielens*. Figures 7a, 7b and 7c demonstrate the performance of the D2P over several values of  $p^*$  on the Movielens dataset. In Figure 7a, we observe that the quality drops only by 3.24%, in terms of Pr@5, when compared to a non-private approach ( $p^* = 1$ ). *Jester.* Figures 8a, 8b and 8c show the results of the performance of the D2P over several values of  $p^*$  on Jester workload. In Figure 8a, we observe that the quality drops only by 2.9% in terms of Pr@5. Interestingly, we observe in Figure 8b that the recall of a non-private approach ( $p^* = 1$ ) is very similar to the one achieved by D2P (e.g, at N = 20, the *recall* values differ by 0.02 only). This observation also means that D2P provides good recommendation quality in datasets with higher rating densities. The higher the profiler probability, the better the recommendation quality.

#### Effect of Selector Probability (p).

Here, we vary the probability p from the *Selector* algorithm from a minimum of 0 to a maximum of 0.5 (with  $\lambda = 1, p^* = 0$ ). *Movielens*. Figures 9a, 9b and 9c demonstrate the performance of D2P over several values of p on Movielens.

*Jester.* Figures 10a, 10b and 10c show the results of the performance of D2P over several values of p on Jester dataset. The lower the selector probability, the better the recommendation quality.

#### *Effect of distance metric* ( $\lambda$ )*.*

We also analyzed the effect of varying the level of privacy using the distance parameter:  $\lambda$ . We observed the quality of recommendations provided by D2P with several values of  $\lambda$  (with p = 0.5,  $p^* = 0$ ). The results of these experiments are given in Figures 12a and 12b. We observe that a lower  $\lambda$  provides better quality because items gets replaced by closer items for lower  $\lambda$ .

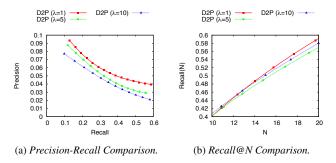


Figure 12: Effect of Distance Metric ( $\lambda$ ) on Quality for the ML Dataset (User-based CF).

#### 4.2.3 Parameter Selection

The distance parameter  $\lambda$  is used to protect user's privacy. We now illustrate its usage on two examples. The first one is depicted in Figure 13. We consider 3 categories (A,B,C), 3 users  $(U_1, U_2, U_3)$  and 5 movies  $(I_1, I_2, I_3, I_4, I_5)$ . We assume that each user wants to hide some specific category. To hide a Category A for user  $U_1$ , we anonymize it with at least one different Category (B or C). We can achieve this by computing the minimum distance for items from Category A in  $U_1$ 's profile  $(I_1, I_3)$  to items in different categories. For item  $I_1$ , we get the distance is 2.8 to  $I_2$  in Category B and 3 to  $I_4$  in Category C. So, the minimum distance for  $I_1$  is 2.8 to  $I_2$  in Category B. We get the same for  $I_3$  in  $U_1$ 's profile. Now, to satisfy the distance for both of these items, we choose the maximum among them which is 2.8. This gives us the  $\lambda_{U_1}$  to hide Category A for  $U_1$ . We do the same for users  $U_2$  and  $U_3$ . Finally, to set the  $\lambda$  for the system, we get the maximum from all users (which is 2.8 in the example).

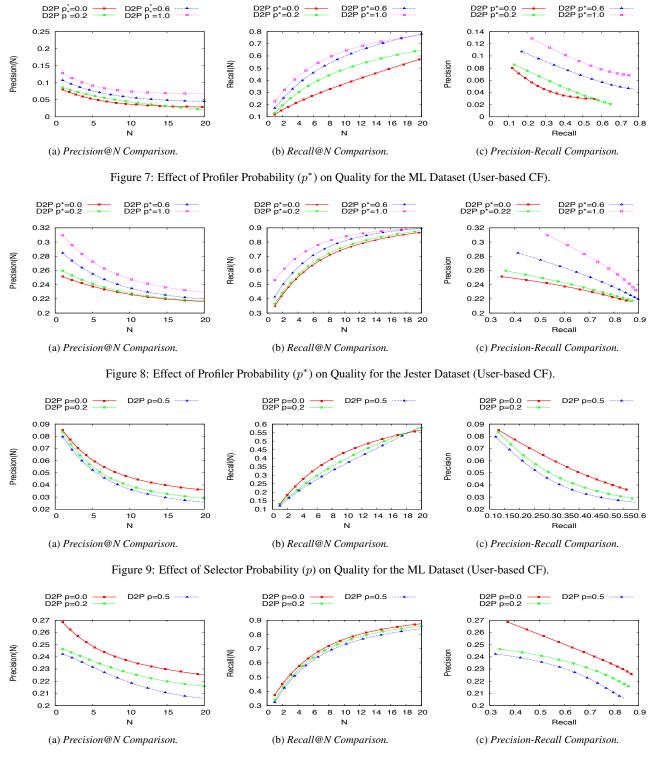


Figure 10: Effect of Selector Probability (p) on Quality for the Jester Dataset (User-based CF).

The distance parameter can be also selected as the average distance for each user profile ( $\lambda_i$ ). Here,  $\lambda_i$  for user  $U_i$  is computed as the average value of the distance between all pairs of items rated by user  $U_i$ . Figure 14 provides an intuition for this distance parameter. For the datasets used for evaluation, we get  $\lambda_{ML1} = 6.5$ ,  $\lambda_{Jester} = 1.5$ . To demonstrate the degradation of  $\epsilon$  based on parameters, p and  $p^*$ , we fix the distance parameter ( $\lambda_{ML1} = 6.5$ ,  $\lambda_{Jester} = 1.5$ ). Figure 15 demonstrates the degradation of  $\epsilon$  based on the privacy parameters (p,  $p^*$ ). For Movielens, we obtain good privacy ( $\epsilon$ = 2.9) and good quality ( $F_1$ -Score=8.5%) with p =0.7,  $p^*$ =0.03,  $\lambda$ =6.5.

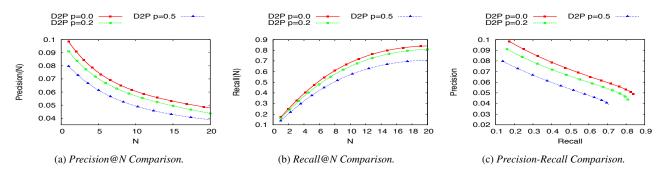


Figure 11: Effect of Selector Probability (p) on Quality for the ML Dataset (Item-based CF).

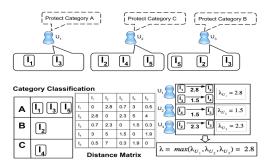


Figure 13: Distance for Personal Choice.

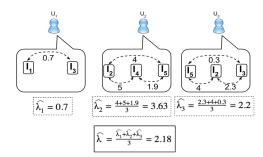


Figure 14: Distance for Average.

For Jester, we obtain good privacy ( $\epsilon$ = 0.97) and good quality ( $F_1$ -Score=23.1%) with p =0.8,  $p^*$ =0.01,  $\lambda$ =1.5.

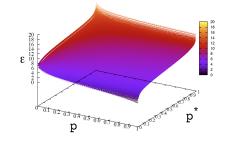
# 4.2.4 Coverage Evaluation

Beyond accuracy, there is a variety of other metrics that should be used to evaluate a recommender [10, 12]. The *Coverage* of a recommender is a metric that captures the domain of items over which it can make recommendations. In particular, we evaluate *Catalog Coverage* [10] of D2P and compare it to the coverage provided by a standard non-private recommender. Consider a set of items  $I_K^j$  contained in the *Top-K* list during the  $j^{th}$  recommendation instance. Also, denote the total number of items by *N*. Hence, *Catalog Coverage* after *M* recommendation instances can be mathematically represented as follows:

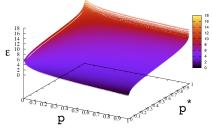
Catalog Coverage = 
$$\frac{|\cup_{j=1...M} I_K^j|}{N}$$

Figure 16 demonstrates the *Catalog Coverage* for D2P and compares it with the coverage in a standard recommender for Movie-

Privacy Parameters for Movielens ( $\lambda$ =6.5)



(a) Privacy Parameters for Movielens (ML1). Privacy Parameters for Jester ( $\lambda$ =1.5)



(b) Privacy Parameters for Jester.

Figure 15: Privacy Parameters Comparison.

lens. We observe that D2P provides 1.5 times better coverage than a standard recommender when the size of recommendation set is 1.

#### 4.2.5 Overhead Evaluation

We evaluate here the computational overhead of D2P's privacy and compare it to the one of [21] which we denote as  $DP_{\delta}$ . We call the computations performed for every recommendation as *Online* computations and the computations done periodically as *Offline* computations. We compare the privacy overhead with the Recommendation Latency (RL) in D2P. Additionally, we compare the privacy overhead in D2P with the privacy overhead in  $DP_{\delta}$ . As shown in Table 4, the overhead for the offline computations in D2P is around 26.4 times smaller than that of [21] for Movielens and around 4.5 times smaller for Jester. All offline computations are parallelised on a 8-core machine.

### 4.2.6 Item-based D2P

D2P can be used with any collaborative filtering technique. We evaluate D2P in another context to illustrate the genericity of D2P.

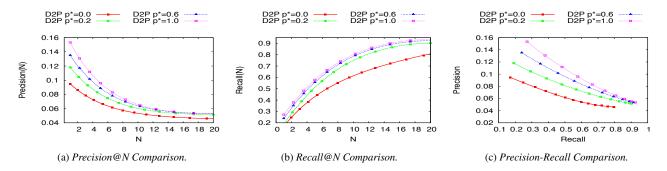


Figure 17: Effect of Profiler Probability  $(p^*)$  on Quality for the ML Dataset (Item-based CF).

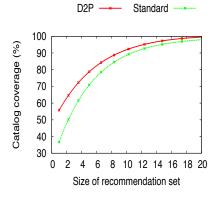


Figure 16: Catalog Coverage Comparison.

Datasets	$\mathcal{D}2\mathcal{P}$ Overhead			$DP_{\delta}$ Overhead
	RL	Online	Offline	Offline
ML1	196ms	32ms	4.54s	120s
Jester	24ms	12ms	162ms	740ms

Table 4: Overhead of Privacy.

We implemented an item-based version of  $D2\mathcal{P}$ :  $i-D2\mathcal{P}$ . In  $i-D2\mathcal{P}$ , the grouping phase is responsible for creating groups of similar *users* based on the distance metric  $\lambda$ . The selector and profiler components in  $i-D2\mathcal{P}$  create *AlterReplica* profiles of the items using the same approach as in  $D2\mathcal{P}$ . Finally, the item recommendations are computed using these *AlterReplica* profiles during the recommendation phase in  $i-D2\mathcal{P}$ . Figures 11a, 11b and 11c convey the quality of recommendations provided by  $i-D2\mathcal{P}$  for varying values of parameter p (with  $\lambda = 1, p^* = 0$ ). Figures 17a, 17b and 17c convey the quality of recommendations provided by  $i-D2\mathcal{P}$  for several values of parameter  $p^*$  (with  $\lambda = 1, p = 0.5$ ). In Figure 17a, we observe that the quality drops by 1.89% in terms of Pr@5 for the ML dataset. This shows that  $D2\mathcal{P}$  also provides good quality of recommendations in item-based CF recommenders.

# 5. RELATED WORK

The notion of differential privacy was introduced by Cynthia Dwork [7, 8, 9]. Most of the research focused on theoretical aspects and provided feasibility and infeasibility results [15]. In this paper, we extend differential privacy to the context of recommenders. We appended the original definition with a distance metric ( $\lambda$ ) and presented an effective way to achieve it through our D2P protocol.

Polat et. al. [25] proposed a randomized perturbation technique to protect user's privacy. Zhang et. al. [34] showed however that a considerable amount of information can be derived from randomly perturbed ratings. Instead of adding perturbations to user profiles, D2P uses the *AlterEgo* profiles which are created based on a distance threshold ( $\lambda$ ). Privacy breaches (compromised user identities) occur when e-commerce sites release their databases to third-parties for data-mining or statistical reporting [2]. The fact that with D2P, the third-parties have only access to the *AlterEgo* profiles alleviates the risk of revealing user's identity to those third parties.

In fact, although, there had been a lot of research work related to privacy in online recommenders [16, 20] and differential privacy [7, 8, 9], only a few of these combined these two notions [13, 21]. McSherry et. al. designed a relaxed version of differential privacy in the context of recommenders [21]. In short, the idea is to add to the ratings - a limited amount of - Gaussian noise. Our notion of distance-based differential privacy provides a stronger form of classical differential privacy in the context of recommender systems. In our case, we replaced items in users profiles with others at some distance. Other differences between the two approaches include the way dynamic updates are addressed as well as the underlying overhead. McSherry et. al. does not consider updates to the covariance matrix, and hence is not applicable to a dynamic system without jeopardizing the privacy guarantee. The AlterEgo profiles used in D2P can grow naturally without the need to recompute from scratch like in [21]. Also, the underlying overhead in D2P is lower. As shown in Table 4, the overhead in D2P is around 26.4 times smaller than that of [21] for Movielens and around 4.5 times smaller for Jester. The additional overhead in [21] stems from the compute-intensive preprocessing steps: (i) removal of per-movie global effects and (ii) centering and clamping process.

# 6. CONCLUDING REMARKS

While personalization has become crucial on the web, it raises however privacy concerns as its quality relies on leveraging user profiles. In this paper, we present an extension of the notion of differential privacy to the context of recommenders: systems that personalize recommendations based on similarities between users. We introduced D2P which ensures this strong form of privacy. D2Paddresses the tradeoff between privacy and quality of recommendation: it can be applied to any collaborative recommender.

The main intuition behind D2P is to rely on a distance metric between items so that groups of similar items can be identified. D2Pleverages this notion of group to generate, from real user profiles, alternative ones, called AlterEgo profiles. These represent differentially private versions of the exact profiles. Such profiles are then used to compute the KNN and provide recommendations. We analyze D2P and evaluate experimentally the impact of the privacy mechanism on the quality of the recommendation in the context of two datasets: MovieLens and Jester. Our results show that privacy can be ensured without significantly impacting the quality of the recommendation. Our experiments demonstrate that D2P can provide 1.5 times better coverage than a standard recommender for Movielens. Additionally, D2P incurs a small privatization overhead compared to other privacy-preserving system like [21] which makes it comparatively more practical for dealing with real-time workloads.

D2P could be further extended to other filtering techniques that rely on user profiles for their recommendation computations. It would also be interesting to incorporate a hybrid approach in D2Pwhere the item groups would be formed using content-based filtering [33] while the actual recommendations would be made based on collaborative filtering techniques. Another interesting direction for future work is to introduce D2P in matrix factorisation techniques [18].

One limitation of D2P stems from the fact that the users trust the service-providers with the original user profiles. Privacy could hence be compromised by online spying on users' activities [11]. For future work, we would like to study the impact on privacy and recommendation quality of probabilistically altering or encrypting user rating [3]: the goal would be to preserve the profile anonymity even from service-providers. Combining such techniques with D2Pwould result in a recommender which is robust to malicious users and even untrusted service-providers engaged in spying activities.

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