# EnviroMeter: A Platform for Querying **Community-Sensed Data**

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# ABSTRACT

Efficiently querying data collected from Large-area Community driven Sensor Networks (LCSNs) is a new and challenging problem. In our previous works, we proposed adaptive techniques for learning models (e.g., statistical, nonparametric, etc.) from such data, considering the fact that LCSN data is typically geo-temporally skewed. In this paper, we present a demonstration of EnviroMeter. EnviroMeter uses our adaptive model creation techniques for processing continuous queries on community-sensed environmental pollution data. Subsequently, it efficiently pushes current pollution updates to GPS-enabled smartphones (through its Android application) or displays it via a web-interface. We experimentally demonstrate that our model-based query processing approach is orders of magnitude efficient than processing the queries over indexed raw data.

#### INTRODUCTION 1.

Community-driven sensing relies on on-board or smartphoneembedded sensors carried by the community (buses, cars, people) to sense an environmental phenomenon of interest (e.g., pollution). The main focus of research until now has been on design and implementation of novel deployments to collect and process community-sensed data. Large-scale community-driven sensor networks are fundamentally different from traditional sensor networks, due to their autonomous and unstructured sensing behavior [5].

An example of such a community-sensed deployment is the OpenSense project [5]. The primary objective of the OpenSense project is to efficiently and effectively monitor environmental pollution using wireless and mobile sensors. The project adopts complex utility driven approaches towards sensing and data management. The geographical granularity for monitoring environmental pollution is on the level of a city or state. Pollution data is collected using sensors installed on public transport buses.

Unfortunately, LCSNs cannot be tightly controlled especially when deployments cover large areas, which makes it Permission to make digital or hard copies of all or part of this work for personal or classroom use is granted without fee provided that copies are not made or distributed for profit or commercial advantage and that copies bear this notice and the full citation on the first page. To copy otherwise, to republish, to post on servers or to redistribute to lists, requires prior specific permission and/or a fee. Articles from this volume were invited to present their results at The 39th International Conference on Very Large Data Bases, August 26th - 30th 2013, Riva del Garda, Trento, Italy. Proceedings of the VLDB Endowment, Vol. 6, No. 12

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difficult to produce a homogeneous view of the phenomenon. Thus, the data collected by sensors is geo-temporally skewed. Data skewness drastically affects the efficiency and accuracy of query processing in the following ways: (a) due to semicontrolled or uncontrolled mobility of sensors, sensor values may not be always available at a particular position and time, (b) the accuracy of an estimated pollution value is not high if a distant sensor value is used to approximate the pollution value at the current position. Queries that are typically processed on such data are of two types: point queries and continuous queries. Point queries return the pollution value at a given position, while continuous queries are registered by a mobile object, which is interested in continuously knowing the pollution around it.

In our prior research [6], we demonstrated that multimodel approaches are suitable for modeling geo-temporally skewed community-sensed data. Previous literature exists on scalable high-speed data processing on the server that addresses efficient storing and querving of models from raw data [9, 10]. In the past, research has focused on the problem of inherent unreliability of the sensors due to their autonomous human-influenced nature or surrounding weather conditions, due to which sensors become error-prone or run out of battery [7, 8].

Although the previous works investigate various ways for the community to perform sensing, they are not concerned with using the collected data and knowledge in providing feedback and information about the phenomenon, back to the community. This work, is focused in closing the gap in the loop. In this demonstration of EnviroMeter, we use our proposed techniques from [6] for succinctly representing community-sensed data in the form of models.

We demonstrate how to use the learned models for efficiently processing user queries. Our techniques adapt to the changing nature of the sensed phenomenon by adjusting the geographical granularity of the models, to capture the phenomena with high fidelity. EnviroMeter is a complete framework that senses data in a community-driven sensor network. It processes the results efficiently, and uses lazy update policies to present them to users on their mobile devices, while significantly reducing network bandwidth and processing delay.

#### 2. SYSTEM DESIGN

The system designed for supporting EnviroMeter consists of three main components. First component uses the adaptive techniques for learning single or multiple models over

the concerned geographical area. Second component consists of the different query processing methods that use the learned models for answering continuous queries. Third component optimizes the bandwidth required for communicating the requests/responses during query processing. In the following sections we discuss each of these components.

## 2.1 Adaptive Models

The architecture of the EnviroMeter framework is shown in Figure 1. It assumes a geographical region  $\mathcal{R}$ , over which environmental pollution is sensed using community-driven approaches. The sensed data is stored in a database in the form of *raw tuples*. The adaptive modeling approach that we propose creates a multi-model abstraction or a *model cover* over the raw tuples dumped in the region  $\mathcal{R}$ . A model cover is defined as a set of models  $\mathcal{M} = \{M_1, \ldots, M_O\}$  that are respectively responsible for modeling the sub-regions  $R_1, R_2, \ldots, R_O$  of  $\mathcal{R}$ . The sub-regions taken together cover the entire region  $\mathcal{R}$ .



Figure 1: Architecture of the framework.

We denote the raw tuple as  $b_i = (t_i, x_i, y_i, s_i)$ , where  $s_i$  is the raw sensor value, and  $t_i$  and  $(x_i, y_i)$  are the time and the position corresponding to the sensor value  $s_i$ . We assume that the model cover is computed using a window of raw tuples  $\mathcal{W}_c = \langle b_i | cH \leq t_i \leq (c+1)H \rangle$ , where c is a positive integer and H is the window length.

We briefly present the adaptive method, called *adaptive* k-means or Ad-KMN, that gave us the best results among many candidates we designed [6]. This method partitions the region  $\mathcal{R}$  adaptively (i.e., only when and where it is necessary) and estimates the models  $M_1, M_2, \ldots, M_O$ . The standard k-means algorithm uses the Euclidean distance for creating the clusters. Instead, in the Ad-KMN method, we use the model approximation error as an additional clustering criteria. An example of the Ad-KMN method on toy data is shown in Figure 2.

Assume that before executing the Ad-KMN method, we compute two centroids  $\mu_1$  and  $\mu_2$  by executing the standard k-means algorithm using the positions  $(x_i, y_i)$  from  $\mathcal{W}_c$  (refer Figure 2(*a*)). Then, (a) we partition the sensor values in  $\mathcal{W}_c$ , such that  $R_1$  and  $R_2$  contain sensor values that are nearest to  $\mu_1$  and  $\mu_2$  respectively, and (b) for the sensor values in  $R_1$  and  $R_2$  we estimate linear regression models  $M_1$  and  $M_2$  and compute the approximation error <sup>1</sup>.

Next, we check whether the approximation error is within a user-defined threshold  $\tau_n$ . In the regions  $(R_1 \text{ or } R_2)$  where the approximation error is greater than  $\tau_n$ , we introduce an additional cluster centroid (equivalent to splitting the region) and re-estimate all the centroids. This procedure is



Figure 2: Example on toy data: (a) initial regions, and (b) two new regions  $R_3$  and  $R_4$  added after an Ad-KMN iteration.

continued until all the regions meet the approximation error threshold  $\tau_n$ . We denote the cluster centroids  $(\mu_1, \ldots, \mu_O)$  as  $\mu$ .

# 2.2 Continuous Query Processing

The query processing framework is depicted in Figure 3. It consists of a mobile object  $v_q$  that transmits the query tuple  $q_l = (t_l, x_l, y_l)$  at time  $t_l$  from position  $(x_l, y_l)$  to the server using mobile data services (GPRS or 3G). Here, we assume a single mobile object (individual or vehicle) continuously querying for pollution around it. The query that we consider is formally defined as follows:

QUERY 1. Continuous Value Query. Given a mobile object  $v_q$  that continuously transmits the query tuple  $q_l = (t_l, x_l, y_l)$  at time  $t_l$ , interpolate the sensor value  $\hat{s}_l$  at position  $(x_l, y_l)$  and transmit it to  $v_q$ .

Here, the sensor value could be any of the pollutants that are typically monitored: carbon dioxide (CO<sub>2</sub>), carbon monoxide (CO), suspended particulate matter, etc. We assume that the mobile object transmits a query tuple with uniform interval, i.e.,  $|t_{l+1} - t_l|$  is always the same. We propose the following three methods for processing Query 1.



Figure 3: Continuous query processing framework.

**Naïve:** In this method, the server does an exhaustive search in the window  $W_c$  to find all the raw tuples that are in a radius r centered at  $(x_l, y_l)$ . Then the interpolated value  $\hat{s}_l$  is computed as the average value of the sensor values  $s_i$ found in the radius r. This interpolated value  $\hat{s}_l$  is then returned to the object  $v_q$ .

Metric Space Indexing: This method is similar to the naïve method, but it uses a metric space index (e.g., R-tree or VP-tree) to enhance the performance of finding the raw tuples in window  $W_c$  that are within radius r of  $(x_i, y_i)$ .

**Model Cover:** This method uses the model cover  $\mathcal{M}$  and the cluster centroids  $\boldsymbol{\mu}$  for query processing. In this method, we first find the cluster centroid  $\mu_*$  in  $\boldsymbol{\mu}$  that is nearest to  $(x_l, y_l)$ . Then the model  $M_* \in (M_1, \ldots, M_O)$  corresponding to  $\mu_*$  is used for interpolating the sensor value  $\hat{s}_l$ .

### 2.3 Bandwidth Optimization Techniques

<sup>&</sup>lt;sup>1</sup>approximation error is the average percentage error compared to the normal range of  $s_i$  in the environment (pollutant specific).

It is a well-known fact that smartphones spend significant amount of battery power and bandwidth in transmitting data via GPRS or 3G data services. In order to optimize the bandwidth- and power-usage, we propose a caching technique referred to as *model-cache*. Model-cache stores the model cover on the smartphone and only queries the server when the cached model cover becomes invalid.

**Model-Cache:** As a system initialization step  $v_q$  sends a model request, denoted as  $e_l$ , to the server (refer Figure 3). In response to  $e_l$  the server sends the following items: (i) the coefficients of all the models in  $\mathcal{M}$ , (ii) the cluster centroids  $\mu$ , and (iii) the time  $t_n$  until which the current model cover is valid.  $v_q$  stores  $(t_n, \mu, \mathcal{M})$  in its local memory.

Now, when the user, who has EnviroMeter running on his/her smartphone, needs a pollution update, a query tuple  $q_l$  is generated. Then, EnviroMeter checks whether  $t_l \leq t_n$ . If  $t_l \leq t_n$ , then it finds the nearest cluster centroid  $\mu_*$  to  $(x_l, y_l)$ . It uses the model  $M_*$  corresponding to  $\mu_*$  for computing the value  $\hat{s}_l$ , without contacting the server. If  $t_l > t_n$ , then the current model cover is invalid, and a new model request  $e_l$  is sent to the server for updating  $(t_n, \boldsymbol{\mu}, \mathcal{M})$ . Since in practice it often happens that mobile objects have limited mobility or predefined trajectories, we save a considerable amount of bandwidth by caching  $(t_n, \boldsymbol{\mu}, \mathcal{M})$ .

In Section 4, we compare the model-cache technique with a baseline technique, which simply responds to each query tuple with the interpolated sensor value  $\hat{s}_l$ , without caching the models. We experimentally demonstrate that modelcache is approximately 50 times bandwidth efficient as compared to the baseline technique.

# 3. DEMONSTRATION

As a part of the demonstration we present the EnviroMeter Android application. The users are presented with a map of Lausanne, Switzerland. EnviroMeter users can quickly find the CO<sub>2</sub> concentration at their current position. The application has the ability to record routes. After a route has been recorded, the user can view it on a map. In addition, the application presents the average pollution level through the route. An informative text indicating whether this value is acceptable according to the OSHA (Occupational Safety and Health Administration) [1] guidelines is displayed. Moreover, the map shows each of the points in the route with a marker whose color varies from green (safe) to red (hazardous CO<sub>2</sub> levels). Finally, users can set up configuration parameters, like the server address and the interval for the position updates using the settings menu.

In the second part, we demonstrate a web interface of EnviroMeter. The web interface can be used in three different modes: continuous query (Query 1), point query and heatmap visualization. In the continuous query mode, users select a set of points that constitute the route, and the application computes dynamically and displays the average CO<sub>2</sub> level for each point on the route. Figure 5(a) shows an example of the web interface, along with the CO<sub>2</sub> concentration at a point clicked by the user. In single point query mode, users click on a point in the map, and the application presents the interpolated CO<sub>2</sub> concentration measured in parts per million (ppm) at that point. Finally, the user can visualize a heatmap of the area Figure 5(b). The emitting points are the centroids computed by the Ad-KMN algorithm with its pollution level. The points are colored in a



Figure 4: (a) Map showing the points in a continuous query, and (b) user-defined settings.

scale going from acceptable (green) to dangerous to human health (red).



(a) query mode (b) heatmap visualization Figure 5: (a) EnviroMeter web interface for single point query, and (b) heatmap of the pollutant concentration.

# 4. EXPERIMENTAL EVALUATION

We perform the experiments on a real dataset collected in the city of Lausanne, Switzerland for the OpenSense project [5]. The dataset is community-sensed by two public transport buses that are equipped with various environmental pollution sensors. For our experiments, we focus only on  $CO_2$ . Our dataset was collected over a period of 1 month and has 176K raw tuples with sampling interval of 60 seconds. We refer to this dataset as *lausanne-data*. In Section 4.1, we compare the various query processing methods for processing Query 1, followed by the experiments on bandwidth optimization methods in Section 4.2.

### 4.1 Query Processing

To evaluate the performance of our query processing methods, we use a varying window size H from 40 to 240 raw tuples (4 hour window), a radius r of 1 km, and error threshold  $\tau_n = 2\%$ . The naïve and the model cover methods are implemented using Python. For testing the metric space indexing methods, we use Python-based implementations of the R-tree [3] and the VP-tree [4]. We use 5000 point queries for comparing the efficiency, accuracy, and memory consumption of all the query processing methods.

Efficiency: Figure 6(a) presents elapsed time for the described scenario. We can observe that the model cover method processes the queries 7.1 times faster as compared

to the VP-tree method for H = 40. In addition, it is 39.4 times faster than the R-tree method for H = 240.

Accuracy: For measuring the accuracy of the obtained results, we compare the naïve method and our proposed model cover method. Recall that the naïve method computes the value  $\hat{s}_l$  as the average of the sensor values that lie in the radius r of the query tuple. Figure 6(b) shows that our method consistently generates a smaller NRMSE (normalized root-mean-square error) than the naïve method. The R-tree and the VP-tree methods are not considered, since they produce the same result as the naïve method.

**Memory Consumption:** For demonstrating the saving in memory due to the model cover method, we use a larger window size H = 5000. We compared the memory required to store: (a) the complete set of points for the naïve method, (b) the index information for the R-tree and VPtree methods, and (c) the models generated by the model cover method. The memory required is accurately measured using the Pympler library [2]. Figure 7(a) presents the average memory required by all the methods averaged over 10 independent runs. Observe that the model cover method dramatically reduces memory consumption and requires approximately 7 times, 70 times and 407 times less memory than the naïve, R-tree and VP-tree methods respectively.



Figure 6: Comparing efficiency and accuracy of query processing. Note the logarithmic scale on the y-axis in (a).

### 4.2 Bandwidth Optimization

To evaluate the bandwidth savings obtained from using the model-cache technique, we use a continuous query of 100 query tuples. We measured the total number of bytes transmitted and received by the mobile device, and the total time to complete the query. From the results presented in Figure 7(b) we can see that the model-cache technique dramatically reduces the memory consumption and query processing time. Compared to the baseline technique (see Section 2.3), model-cache requires 113 times less transmitted bytes, 30 times less received bytes, and approximately 100 times less time.

### 5. CONCLUSION

In this demonstration, we introduced EnviroMeter, an efficient and easy-to-use framework for processing queries over community-sensed data. We proposed various design strategies for query processing and optimizing network bandwidth. We presented the main functionalities of the EnviroMeter Android application for smartphones and the web interface. Finally, we evaluated our techniques on a real dataset and



Figure 7: (a) Comparing memory requirements of various query processing methods and (b) comparing the bandwidth optimization techniques. Note the logarithmic scale on the y-axis.

clearly demonstrated the orders of magnitude performance enhancements obtained using our methods.

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